

Pessimistic: Exclusive Locking

- Each object has a lock
- At most one transaction can be inside lock
- Before reading or writing object O, transaction T must call `lock(O)`
 - Blocks if another transaction already inside lock
- After entering lock T can read and write O multiple times
- When done (or at commit point), T calls `unlock(O)`
 - If other transactions waiting at `lock(O)`, allows one of them in
- Sound familiar? (This is Mutual Exclusion!)