Views

- Each process maintains a membership list
- The membership list is called a *View*
- An update to the membership list is called a *View Change*
  - Process join, leave, or failure
- Virtual synchrony guarantees that all **view changes are delivered in the same order at all correct processes**
  - If a correct P1 process receives views, say {P1}, {P1, P2, P3}, {P1, P2}, {P1, P2, P4} then
  - Any other correct process receives the *same sequence* of view changes (after it joins the group)
    - P2 receives views {P1, P2, P3}, {P1, P2}, {P1, P2, P4}
- Views may be delivered at different <u>physical</u> times at processes,   but they are delivered in the same <u>order</u>